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# Conjugacy and Dynamics in Thompson's Groups

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**Abstract** We give a unified solution to the conjugacy problem for Thompson's groups F, T, and V. The solution uses "strand diagrams", which are similar in spirit to braids and generalize tree-pair diagrams for elements of Thompson's groups. Strand diagrams are closely related to piecewise-linear functions for elements of Thompson's groups, and we use this correspondence to investigate the dynamics of elements of F. Though many of the results in this paper are known, our approach is new, and it yields elegant proofs of several old results.

#### 1 Introduction

In the late 1960's, Richard J. Thompson introduced three groups F, T, and V, now known collectively as **Thompson's groups**. The group F consists of all piecewise-linear homeomorphisms of [0,1] with finitely many breakpoints satisfying the following conditions:

- 1. Every slope is a power of two, and
- 2. Every breakpoint has dyadic rational coordinates.

The groups T and V are defined similarly, except that T is a group of homeomorphisms of the circle, and V is a group of homeomorphisms of the Cantor set. These three groups have been studied extensively, and we will assume that they are somewhat familiar to the reader. For a thorough introduction, see the notes by Cannon, Floyd and Parry [7].

In this paper, we present a unified solution to the conjugacy problems for Thompson's groups F, T, and V. The conjugacy problems for F and V have been solved separately

The first author gratefully acknowledges partial support from an NSF Postdoctoral Research Fellowship while he was at Texas A&M University.

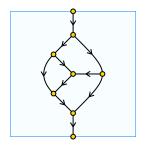
This work is part of the second author's Ph.D. thesis at Cornell University. The second author gratefully acknowledges the Centre de Recerca Matemàtica (CRM) and its staff for the support received during the development of this work.

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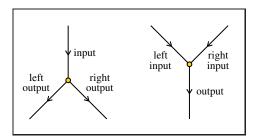


Fig. 2.1: A strand diagram.

Fig. 2.2: A split and a merge.

before: a solution for F was given by Guba and Sapir [10] in the more general context of diagram groups, while a solution for V was given by Higman [11] and again by Salazar-Diaz [21]. As far as we know, the solution to the conjugacy problem for Thompson's group T is new. In addition, our methods work the same way for all three groups, and may be helpful for solving the conjugacy problem for certain related groups, such as the generalized Thompson groups F(p) defined by Brown [6], the universal central extension  $\widetilde{T}$  of T [1], or the braided Thompson group BV defined by Brin [4].

Our solution uses **strand diagrams** for elements of Thompson's groups. These are a generalization of the standard tree-pair diagrams, and were used in [1] to construct an Eilenberg-MacLane space for Thompson's group F. Roughly speaking, a strand diagram is similar to a braid, but instead of twists a strand diagram has splits and merges. In the case of F, strand diagrams are closely related to the "diagrams" of Guba and Sapir [10,9], and are essentially equivalent to the "pictures" of Bogley and Pride [3,19,20].

In addition to solving the conjugacy problem, our invariants yield information about the dynamics of elements of Thompson's groups. This arises from an explicit correspondence between strand diagrams and piecewise-linear functions, which we describe in Section 5. Using this correspondence, we obtain a complete understanding of the dynamics of elements of F, giving simple proofs of several previously known results. A more thorough investigation of dynamics using strand diagrams can be found in [15].

# 2 Strand Diagrams

In this section we introduce strand diagrams, which will be our main tool for solving the conjugacy problem. We begin by realizing F, T, and V as groups of strand diagrams, and then continue by defining the annular, toral, and closed strand diagrams that will serve as our conjugacy invariants.

#### 2.1 Strand Diagrams for F, T, and V

**Definition 2.1** A **strand diagram** (see Figure 2.1) is a finite acyclic digraph embedded in the unit square  $[0,1] \times [0,1]$ , with the following properties:

- 1. The graph has a single univalent source, which lies on the top edge of the square, and a single univalent sink, which lies on the bottom edge.
- 2. Every other vertex is trivalent, and is either a **split** or a **merge**, as shown in Figure 2.2.

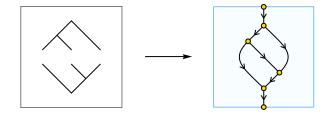


Fig. 2.3: Constructing a strand diagram from a tree-pair diagram.

By convention, isotopic strand diagrams are considered equal.

Strand diagrams are a generalization of tree-pair diagrams for elements of Thompson's group F. Specifically, given any tree-pair diagram, we can construct a corresponding strand diagram by gluing together the leaves of the two trees, as shown in Figure 2.3. Not every strand diagram can be obtained in this fashion, so we can view tree-pair diagrams as a proper subset of strand diagrams.

**Definition 2.2** A **reduction** of a strand diagram is either of the two moves shown in Figure 2.4. A strand diagram is **reduced** if it is not subject to any reductions. Two strand diagrams are **equivalent** if one can be obtained from the other by a sequence of reductions and inverse reductions.

Note that two tree-pair diagrams for the same element of F yield equivalent strand diagrams. In particular, we can use type I reductions to eliminate redundant pairs of carets that have been glued together.

**Proposition 2.3** Every strand diagram is equivalent to a unique reduced strand diagram.

*Proof* This is a straightforward application of the theory of abstract rewriting systems. In particular, the process of reduction is terminating since each reduction decreases the number of vertices, and it is straightforward to check that reductions are locally confluent.

Every reduced strand diagram can be obtained by gluing together the trees of a reduced tree-pair diagram; thus, this proposition gives us a one-to-one correspondence between equivalence classes of strand diagrams and elements of Thompson's group F. We will make this correspondence more explicit in Section 5 when we investigate the relationship between strand diagrams and piecewise-linear functions.

The advantage of strand diagrams over tree-pair diagrams is that strand diagrams are easier to compose:

<sup>&</sup>lt;sup>1</sup> A weaker version of the Newman's Diamond Lemma suffices: see Theorem 1 in [18].



Fig. 2.4: The two reductions.

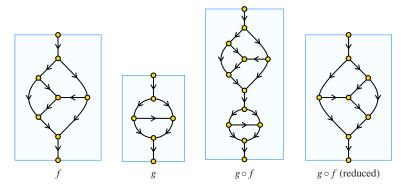


Fig. 2.5: Composing elements using strand diagrams.

**Definition 2.4** The **concatenation** of two strand diagrams is obtained by gluing the sink of the first to the source of the second, and then eliminating the resulting bivalent vertex (see Figure 2.5).

Concatenation of strand diagrams corresponds to composition of elements of F, as shown in Figure 2.5. Note that the result is usually not reduced, so in practice we compose two elements by concatenating the strand diagrams and then reducing the result.

The following proposition should be clear from the above discussion:

**Proposition 2.5** The set of all reduced strand diagrams forms a group under the operation of concatenation followed by reduction. This is isomorphic to Thompson's group F.

For technical reasons, we will also need to consider strand diagrams with more than one source or sink:

**Definition 2.6** Let m and n be positive integers. An (m, n)-strand diagram is similar to a strand diagram, except that it has m univalent sources along the top of the square, and n univalent sinks along the bottom of the square (see Figure 2.6).

Note that a "strand diagram" is the same thing as a (1,1)-strand diagram. As with strand diagrams, every (m,n)-strand diagram is equivalent to a unique reduced (m,n)-strand diagram. Moreover, every reduced (m,n)-strand diagram can be obtained by attaching together the leaves of two binary forests, one with m trees and the other with n trees.

The set of all reduced (m,n)-strand diagrams forms a groupoid over the positive integers. In particular, we can concatenate any (i,j)-strand diagram with any (j,k)-strand diagram

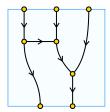
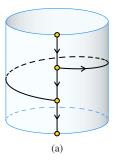
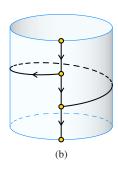


Fig. 2.6: A (3,2)-strand diagram.





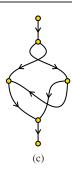


Fig. 2.7: (a) A cylindrical strand diagram. (b) An isotopic cylindrical strand diagram, obtained from the first by a Dehn twist. (c) An abstract strand diagram.

by attaching the sinks of the first to the sources of the second. The isotropy group of this groupoid at the point 1 (or indeed at any point) is isomorphic to Thompson's group F.

We now turn to strand diagrams for Thompson's groups T and V. For Thompson's group T, we can use the following class of diagrams:

**Definition 2.7** A **cylindrical strand diagram** (see Figure 2.7a) is a finite acyclic digraph embedded in the cylinder  $S^1 \times [0, 1]$ , with the following properties:

- 1. The graph has a single univalent source, which lies on the top circle of the cylinder, and a single univalent sink, which lies on the bottom circle of the cylinder.
- 2. Every other vertex is either a split or a merge.

Isotopic cylindrical strand diagrams are considered equal.

Note that the source and sink of a cylindrical strand diagram may rotate around their respective circles during an isotopy. Therefore, two cylindrical strand diagrams that differ by a Dehn twist of the cylinder are considered equal (see Figure 2.7b).

Note also that we can glue two binary trees together on the cylinder via any cyclic permutation of the leaves. This lets us view tree-pair diagrams for elements of T as a proper subset of cylindrical strand diagrams.

It is less obvious how to define strand diagrams for Thompson's group V. For the following definition, recall that a **rotation system** on a graph is an assignment of a circular order to the edges incident on each vertex [17].

**Definition 2.8** An **abstract strand diagram** is a finite acyclic digraph together with a rotation system, having the following properties:

- 1. The graph has a single univalent source and a single univalent sink.
- 2. Every other vertex is either a split or a merge.

Two abstract strand diagrams are considered equal if there exists a digraph isomorphism between them that is compatible with the corresponding rotation systems.

Figure 2.7c shows an abstract strand diagram drawn in the plane. By convention, the rotation system is indicated by the counterclockwise order of the edges at each vertex. Note that the rotation system on an abstract strand diagram lets us define the left and right outputs for a split, and the left and right inputs for a merge.



Fig. 2.8: Invalid reductions.

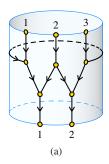
Given a tree-pair diagram for an element of V, we can glue the leaves of the trees together along the indicated bijection to obtain an abstract strand diagram. This lets us view tree pair diagrams for elements of V as a subset of abstract strand diagrams.

All of the constructions we defined for strand diagrams can be extended to cylindrical and abstract strand diagrams. In particular:

1. We can reduce either type of diagram using the two reduction moves shown in Figure 2.4. When applying these moves, it is very important to respect the counterclockwise orientation of the edges shown in Figure 2.4. In particular, neither of the moves shown in Figure 2.8 is a valid reduction.

Assuming we apply only valid reduction moves, every cylindrical strand diagram is equivalent to a unique reduced cylindrical strand diagram, and every abstract strand diagram is equivalent to a unique reduced abstract strand diagram.

- 2. We can concatenate two cylindrical strand diagrams or two abstract strand diagrams. Under the operation of concatenation followed by reduction, reduced cylindrical strand diagrams form a group isomorphic to Thompson's group *T*, and reduced abstract strand diagrams form a group isomorphic to Thompson's group *V*.
- 3. We can also define groupoids of cylindrical (m,n)-strand diagrams and abstract (m,n)-strand diagrams, though there is a slight complication. If we wish to be able to concatenate (m,n)-diagrams unambiguously, we must include numberings of the sources and sinks, as shown in Figure 2.9. For a cylindrical diagram, these numbers are required to appear in counterclockwise order around each circle. For an abstract diagram, the sources and sinks can be numbered in any order. When concatenating two diagrams, we use these numbers to determine which sources should be attached to which sinks.



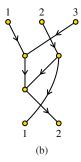


Fig. 2.9: (a) A cylindrical (3,2)-strand diagram. (b) An abstract (3,2)-strand diagram.

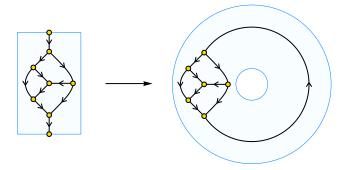


Fig. 2.10: Closing a strand diagram to obtain an annular strand diagram.

#### 2.2 Conjugacy Invariants

Given a strand diagram in the unit square, we can **close** it to obtain a graph embedded in an annulus, as shown in Figure 2.10. More generally, we can close any (m, n)-strand diagram for which m = n. This prompts the following definition:

**Definition 2.9** An **annular strand diagram** is a finite directed topological graph embedded in the annulus  $[0, 1] \times S^1$ , with the following properties:

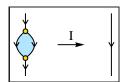
- 1. Every vertex is either a split or a merge.
- 2. Every directed cycle winds counterclockwise around the central hole.

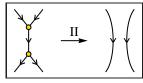
As with strand diagrams, isotopic annular strand diagrams are considered equal.

In this definition, a **topological graph** refers to a graph without bivalent vertices, which may also contain one or more **free loops** (closed loops without any vertices). Note that a free loop counts as a directed cycle, so by property (2) every free loop in an annular strand diagram must wind counterclockwise around the central hole.

Though the underlying spaces are homeomorphic, an annular strand diagram is actually very different from a cylindrical strand diagram. In a cylindrical strand diagram, the horizontal direction is a circle, but there are still sources and sinks along the top and bottom. In an annular strand diagram, it is the vertical direction that has been made into a circle. We shall consistently distinguish between these cases by the order of the factors in a product:  $S^1 \times [0,1]$  for the cylinder, and  $[0,1] \times S^1$  for the annulus.

We can reduce annular strand diagrams using the three reduction moves shown in Figure 2.11. The third move is necessary to make reductions of annular strand diagrams locally





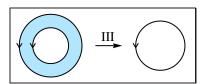


Fig. 2.11: The three reduction moves for an annular strand diagram. In the first move, the shaded region must be a topological disk. In the third move, both loops must be free loops, and the region between must be a topological annulus that does not contain any vertices.

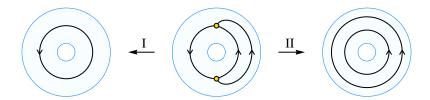


Fig. 2.12: The annular strand diagram in the center can be reduced using either a type I move or a type II move. A type III move is required to reconcile the results.

confluent. In particular, Figure 2.12 shows an annular strand diagram which can be reduced using either a type I move or a type II move, but for which the two results can only be reconciled using a type III move. As long as we use all three reductions, though, every annular strand diagram is equivalent to a unique reduced annular strand diagram.

The following theorem explains our interest in annular strand diagrams. We will prove this theorem in Section 3.

**Theorem 2.10** Let f and g be elements of Thompson's group F. Let f and g be strand diagrams for f and g, and let f' and g' be the reduced annular strand diagrams obtained by closing f and g and reducing. Then f and g are conjugate if and only if f' and g' are isotopic.

This theorem provides a geometric solution to the conjugacy problem in Thompson's group F. (Recall that isotopy of graph embeddings can be checked algorithmically using dual graphs.) This procedure is surprisingly fast: it is shown in [12] that an algorithm based on the above theorem can be implemented in linear time.

Theorem 2.10 is very similar to the solution to the conjugacy problem given by Guba and Sapir in [10], using "diagrams" instead of strand diagrams. Indeed, our strand diagrams are essentially just the dual graphs to the diagrams that they use. The advantage of strand diagrams is that they generalize easily to Thompson's groups T and V.

We begin by describing the generalization to Thompson's group T. Given a cylindrical strand diagram, we can **close** it by gluing together the two circles of the cylinder, identifying the source and the sink (see Figure 2.13). The result is a graph embedded on a torus  $S^1 \times S^1$ .

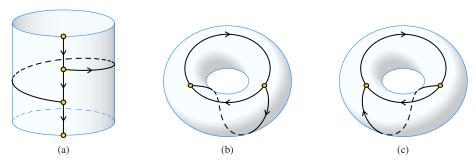


Fig. 2.13: (a) A cylindrical strand diagram. (b) The corresponding toral strand diagram. (c) The same toral strand diagram after a Dehn twist around  $S^1 \times \{1\}$ .

For the following definition, let c denote the cohomology class in  $H^1(S^1 \times S^1)$  that measures the intersection number with the circle  $S^1 \times \{1\}$  (the circle obtained by gluing together the top and bottom circles of the cylinder). Note that this cohomology class is invariant under the Dehn twist around the circle  $S^1 \times \{1\}$ , and indeed the stabilizer of c in the mapping class group is precisely the cyclic subgroup generated by this Dehn twist.

**Definition 2.11** A **toral strand diagram** is a finite directed topological graph embedded on the torus  $S^1 \times S^1$ , with the following properties:

- 1. Every vertex is either a split or a merge.
- 2. Every directed cycle evaluates to a positive value under c.

Two toral strand diagrams are considered equal if they are isotopic, or if they differ by finitely many Dehn twists around the circle  $S^1 \times \{1\}$  (see Figure 2.13c).

We can reduce toral strand diagrams in exactly the same way that we reduce annular strand diagrams, and each toral strand diagram is equivalent to a unique reduced toral strand diagram. The following theorem provides a solution to the conjugacy problem in Thompson's group *T*. We will prove this theorem in Section 3.

**Theorem 2.12** Let f and g be elements of Thompson's group T. Let f and g be cylindrical strand diagrams for f and g, and let f' and g' be the reduced toral strand diagrams obtained by closing f and g and reducing. Then f and g are conjugate if and only if f' and g' are equal.

Note that checking whether  $\mathfrak{f}'$  and  $\mathfrak{g}'$  are equal involves checking whether they differ by finitely many Dehn twists around  $S^1 \times \{1\}$ . This can be determined algorithmically as follows:

- 1. Enumerate all digraph isomorphisms between the two diagrams.
- 2. For each digraph isomorphism, use dual graphs to check whether it extends to a self-homeomorphism of the torus.
- 3. For each digraph isomorphism that extends to a self-homeomorphism of the torus, check whether this self-homeomorphism preserves the cohomology class *c*. If it does, then the two toral strand diagrams are equal. If no such isomorphism is found, then the two toral strand diagrams are different.

Here c is the conjugacy class that measures winding number around the central hole.

Finally, we would like to discuss the solution to the conjugacy problem in Thompson's group V. Because abstract strand diagrams for V do not lie on a surface, we must replace the "central hole" of the annulus or torus with a cohomology class:

**Definition 2.13** An **abstract closed strand diagram** is an ordered pair  $(\Gamma, c)$ , where  $\Gamma$  is a finite directed topological graph with a rotation system, and  $c \in H^1(\Gamma; \mathbb{Z})$ , having the following properties:

- 1. Every vertex of  $\Gamma$  is either a split or a merge.
- 2. For every directed cycle  $\alpha$  in  $\Gamma$ , we have  $c(\alpha) > 0$ .

Two abstract closed strand diagrams  $(\Gamma, c)$  and  $(\Gamma', c')$  are considered equal if there exists an isomorphism  $\varphi \colon \Gamma \to \Gamma'$  of directed graphs with rotation systems such that  $\varphi^*(c') = c$ .

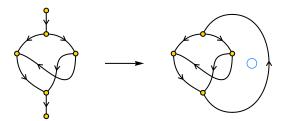


Fig. 2.14: Closing an abstract strand diagram.

We can **close** any abstract strand diagram by attaching the source and the sink, resulting in an abstract closed strand diagram (see Figure 2.14). The cohomology class c is defined by the 1-cochain that evaluates to 1 on the new edge obtained by gluing the source to the sink, and 0 on every other edge.

Figures 2.14 and 2.15 show drawings of several abstract closed strand diagrams. In each drawing, the rotation system r is indicated by the counterclockwise order of the edges around each vertex, and the cohomology class c corresponds to counterclockwise winding number around the indicated "hole". It is always possible to draw an abstract closed strand diagram in this way, since the elements of  $H^1(\Gamma;\mathbb{Z})$  are in one-to-one correspondence with homotopy classes of maps from  $\Gamma$  to the punctured plane.

We shall refer to the cohomology class c as the **cutting class** of the abstract closed strand diagram. Note that the cutting class is an essential part of the invariant, and is not determined by the underlying graph  $\Gamma$ . For example, Figure 2.15 shows two abstract closed strand diagrams with isomorphic graphs and rotation systems, but slightly different cutting classes. These two diagrams are not equal, and correspond to different conjugacy classes in Thompson's group V.

*Remark 2.14* Though we can reduce an abstract closed strand diagram using the three reductions shown in Figure 2.11, some care must be taken in interpreting these moves. In particular, the following restrictions apply:

- 1. For a type I reduction, the closed cycle  $\alpha$  formed by the two parallel edges must satisfy  $c(\alpha) = 0$ . This replaces the requirement that the two parallel edges bound a disk.
- 2. For a type III reduction, the two free loops  $\beta$  and  $\gamma$  must satisfy  $c(\beta) = c(\gamma)$ . This replaces the requirement that the two free loops bound an annulus.

In addition, the counterclockwise order of the edges at each vertex must be exactly as shown in Figure 2.11. That is, the two moves shown in Figure 2.8 are not allowed.

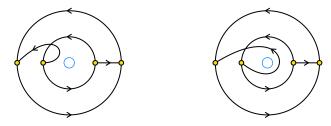


Fig. 2.15: Two abstract closed strand diagrams with slightly different cutting classes.

Note also that, for each type of reduction, the cutting class on the original diagram induces a cutting class on the reduced diagram in a natural way.

The following theorem characterizes conjugacy in Thompson's group V. We will prove it in Section 3.

**Theorem 2.15** Let f and g be elements of Thompson's group V. Let f and g be abstract strand diagrams for f and g, and let f' and g' be the reduced abstract closed strand diagrams obtained by closing f and g and reducing. Then f and g are conjugate if and only if f' and g' are equal.

Note that we can check equality of abstract closed strand diagrams algorithmically by enumerating all isomorphisms between the directed graphs and then checking whether the isomorphisms are compatible with the rotation systems and cutting classes. Thus the above theorem provides a solution to the conjugacy problem in V.

#### 3 Solving the Conjugacy Problem

The goal of this section is to prove Theorems 2.10, 2.12, and 2.15, which solve the conjugacy problems for F, T, and V. Our approach is to handle all three cases simultaneously, pointing out the differences when necessary.

Throughout this section, we will use the term (m,n)-diagram to refer to either an (m,n)-strand diagram, a cylindrical (m,n)-strand diagram, or an abstract (m,n)-strand diagram. Similarly, we will use the term **closed diagram** to refer to either an annular strand diagram, a toral strand diagram, or an abstract closed strand diagram.

Recall that two diagrams are **equivalent** if they differ by a sequence of reductions and inverse reductions. For convenience we will be working primarily in the groupoid of equivalence classes of (m,n)-diagrams. If  $\mathfrak{f}$  is an (m,n)-diagram, we will let  $[\mathfrak{f}]$  denote the corresponding element of the groupoid. Here is our main theorem:

**Theorem 3.1** Let  $\mathfrak{f}$  be an (m,m)-diagram, and let  $\mathfrak{g}$  be an (n,n)-diagram. Then  $[\mathfrak{f}]$  and  $[\mathfrak{g}]$  are conjugate if and only if the closures of  $\mathfrak{f}$  and  $\mathfrak{g}$  are equivalent.

Note that Theorems 2.10, 2.12, and 2.15 follow immediately. In particular, two elements of F, T, and V are conjugate within the group if and only if they are conjugate within the corresponding groupoid.

The proof of Theorem 3.1 occupies the remainder of this section. Most of the content of the proof is covered in the following pair of propositions.

**Proposition 3.2** Let  $\mathfrak{f}$  be an (m,m)-diagram, let  $\mathfrak{f}'$  be its closure, and let  $\mathfrak{g}'$  be a closed diagram obtained by applying a reduction to  $\mathfrak{f}'$ . Then there exists an (n,n)-diagram  $\mathfrak{g}$  whose closure is  $\mathfrak{g}'$  such that  $[\mathfrak{g}]$  is conjugate to  $[\mathfrak{f}]$ .

*Proof* Throughout this proof, we will use the term "gluing points" to refer to the points in  $\mathfrak{f}'$  obtained by gluing together the sources and sinks of  $\mathfrak{f}$ . There are three cases, based on the type of reduction.

Suppose the reduction of  $\mathfrak{f}'$  is of type I. If this corresponds to a type I reduction on  $\mathfrak{f}$ , then we are done. However, it is possible that the reduction has a gluing point on each of the parallel edges, as shown in Figure 3.1a. In this case, conjugating by an element that merges the two sinks will produce  $\mathfrak{g}$ , as shown in the figure. Finally, it is possible that each of the

parallel edges contains several gluing points, as shown in Figure 3.1b. In this case, we must conjugate by an element that merges each pair of relevant sinks. Note that the number of gluing points must be the same on the two parallel edges because of the requirement that they bound a disk (see Figure 2.11)—or, in the case of abstract diagrams, because of the cohomology requirement stated in Remark 2.14.

The argument for a type II reduction is fairly similar. If the reduction corresponds to a reduction on f, we are done. Otherwise we must conjugate by an element with one or more splits, as shown in Figures 3.1c and Figures 3.2a.

Finally, suppose the reduction of  $\mathfrak{f}'$  is of type III. If there is one gluing point on each free loop, then we can conjugate by an element with a single split, as shown in Figure 3.2b. This is the only possibility in the annular case, but in the toral and abstract cases we might have multiple gluing points on each free loop, as shown in Figure 3.2c. In this case, we must conjugate by an element with multiple splits, as shown. Again, note that the number of gluing points must be the same on each loop because of the requirement that they bound an annulus (see Figure 2.11)—or, in the case of abstract diagrams, because of the cohomology requirement stated in Remark 2.14.

**Proposition 3.3** Let  $\mathfrak{f}$  be an (m,m)-diagram, let  $\mathfrak{g}$  be an (n,n)-diagram. If the closures of  $\mathfrak{f}$  and  $\mathfrak{g}$  are equal, then  $[\mathfrak{f}]$  and  $[\mathfrak{g}]$  are conjugate.

*Proof* Let  $\{f_k\}_{k\in\mathbb{Z}}$  be a collection of disjoint copies of  $\mathfrak{f}$ , and consider the "infinite strand diagram"  $\mathfrak{f}^{\infty}$  obtained by gluing the sinks of each  $\mathfrak{f}_k$  to the sources of  $\mathfrak{f}_{k+1}$ . If  $\mathfrak{f}$  is annular or toral, then  $\mathfrak{f}^{\infty}$  is an infinite graph embedded in  $[0,1] \times \mathbb{R}$  or  $S^1 \times \mathbb{R}$ , respectively. If  $\mathfrak{f}$  is abstract, then  $\mathfrak{f}^{\infty}$  is simply an infinite directed graph with a rotation system.

Let  $\mathfrak{f}'$  be the closure of  $\mathfrak{f}$ , and note that  $\mathfrak{f}^{\infty}$  is an infinite-sheeted cover of  $\mathfrak{f}'$ . In the annular or toral case, this cover is induced by the covering map from  $[0,1]\times\mathbb{R}$  to the annulus, or from  $S^1\times\mathbb{R}$  to the torus. In the abstract case, the kernel of the cutting class of  $\mathfrak{f}'$  is an infinite-index subgroup of  $\pi_1(\mathfrak{f}')$ , and  $\mathfrak{f}^{\infty}$  is the cover corresponding to this subgroup. Either way, the cover is entirely determined by  $\mathfrak{f}'$ .

Since the closure  $\mathfrak{g}'$  of  $\mathfrak{g}$  is equal to  $\mathfrak{f}'$ , it follows that the infinite strand diagram  $\mathfrak{g}^{\infty}$  must be the same as  $\mathfrak{f}^{\infty}$ . To be precise:

1. If f' and g' are annular, then they are isotopic, and this isotopy lifts to an isotopy between  $f^{\infty}$  and  $g^{\infty}$ .

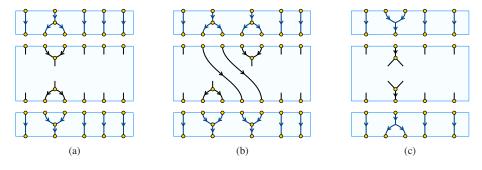
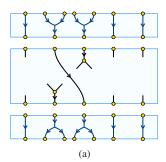
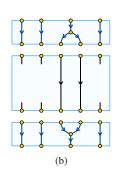


Fig. 3.1: (a) Conjugating  $\mathfrak f$  to perform a type I reduction on  $\mathfrak f'$ . (b) A type I reduction involving more than one pair of gluing points. (c) A type II reduction.





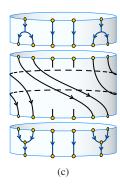


Fig. 3.2: (a) A type II reduction involving more than one gluing point. (b) A type III reduction. (c) A type III reduction involving more than one pair of gluing points.

- 2. If f' and g' are toral, then either they are isotopic, or they differ by finitely many Dehn twists around S¹ × {1}. An isotopy on the torus lifts to an isotopy on the infinite cylinder, and a Dehn twist around S¹ × {1} also lifts to an isotopy on the infinite cylinder. Either way, f<sup>∞</sup> and g<sup>∞</sup> are isotopic.
- 3. If  $\mathfrak{f}'$  and  $\mathfrak{g}'$  are abstract, then there is an isomorphism between them that preserves the cutting class, and this lifts to an isomorphism between  $\mathfrak{f}^{\infty}$  and  $\mathfrak{g}^{\infty}$ .

Therefore,  $\mathfrak{f}^{\infty}$  can be expressed either as the union  $\bigcup_{k\in\mathbb{Z}}\mathfrak{f}_k$  of infinitely many copies of  $\mathfrak{f}$ , or as the union  $\bigcup_{k\in\mathbb{Z}}\mathfrak{g}_k$  of infinitely many copies of  $\mathfrak{g}$ . Moreover, there exists a deck transformation t of  $\mathfrak{f}^{\infty}$  such that  $t(\mathfrak{f}_k) = \mathfrak{f}_{k+1}$  and  $t(\mathfrak{g}_k) = \mathfrak{g}_{k+1}$  for each k.

Now, since  $\mathfrak{f}$  and  $\mathfrak{g}$  are compact, there exists a sufficiently large  $N \in \mathbb{N}$  so that  $\mathfrak{g}_N$  is entirely contained in  $\bigcup_{k>0} \mathfrak{f}_k$ . Let

$$\mathfrak{h}_0 \, = \, \left( \bigcup_{k \geq 0} \mathfrak{f}_k \right) \cap \left( \bigcup_{k < N} \mathfrak{g}_k \right) \qquad \text{and} \qquad \mathfrak{h}_1 \, = \, \left( \bigcup_{k > 0} \mathfrak{f}_k \right) \cap \left( \bigcup_{k \leq N} \mathfrak{g}_k \right).$$

Then  $\mathfrak{h}_0$  and  $\mathfrak{h}_1$  are strand diagrams and  $\mathfrak{h}_1 = t(\mathfrak{h}_0)$ , so  $\mathfrak{h}_0$  and  $\mathfrak{h}_1$  are equal to the same strand diagram  $\mathfrak{h}$ . But  $\mathfrak{f}_0 \cup \mathfrak{h}_1 = \mathfrak{h}_0 \cup \mathfrak{g}_N$ , so the concatenation of  $\mathfrak{f}$  with  $\mathfrak{h}$  is equal to the concatenation of  $\mathfrak{h}$  with  $\mathfrak{g}$ .

We are now ready to prove the main theorem.

*Proof* (*Theorem 3.1*) Let  $\mathfrak{f}$  be an (m,m)-diagram, and let  $\mathfrak{g}$  be an (n,n)-diagram of the same type. Suppose first that  $[\mathfrak{f}]$  and  $[\mathfrak{g}]$  are conjugate. Then  $\mathfrak{f}$  is equivalent to a concatenation of the form  $\mathfrak{h}^{-1}\mathfrak{g}\mathfrak{h}$  for some (m,n)-diagram  $\mathfrak{h}$ . Then the closure of  $\mathfrak{f}$  is equivalent to the closure of  $\mathfrak{h}^{-1}\mathfrak{g}\mathfrak{h}$ . This is the same as the closure of  $\mathfrak{g}\mathfrak{h}\mathfrak{h}^{-1}$ , which is equivalent to the closure of  $\mathfrak{g}$ .

For the converse, suppose that closures of  $\mathfrak{f}$  and  $\mathfrak{g}$  are equivalent. Then applying reductions to either  $\mathfrak{f}$  or  $\mathfrak{g}$  eventually results in the same reduced closed diagram  $\mathfrak{h}'$ . By Proposition 3.2, there exist diagrams  $\mathfrak{h}_1$  and  $\mathfrak{h}_2$  whose closures are  $\mathfrak{h}'$  such that  $[\mathfrak{f}]$  is conjugate to  $[\mathfrak{h}_1]$ , and  $[\mathfrak{g}]$  is conjugate to  $[\mathfrak{h}_2]$ . By Proposition 3.3, the elements  $[\mathfrak{h}_1]$  and  $[\mathfrak{h}_2]$  must themselves be conjugate, and therefore  $[\mathfrak{f}]$  is conjugate to  $[\mathfrak{g}]$ .

### 4 Structure of Reduced Closed Strand Diagrams

In this section we briefly investigate the structure of reduced closed strand diagrams. While reading the subsequent proof, it might help to refer to Figure 4.1, which gives two examples of reduced annular strand diagrams.

As in the previous section, we will used the phrase **closed diagram** to refer to either an annular strand diagram, a toral strand diagram, or an abstract closed strand diagram. A closed strand diagram is **reduced** if it is not subject to any of the three reduction moves shown in Figure 2.11.

We will use the following terminology for certain kinds of directed cycles:

- 1. A **free loop** is a directed cycle with no vertices.
- 2. A **split loop** is a directed cycle with splits, but no merges.
- 3. A **merge loop** is a directed cycle with merges, but no splits.

For example, the first annular strand diagram in Figure 4.1 has one split loop and two merge loops, while the second annular strand diagram in the figure has two split loops, two merge loops, and one free loop.

#### **Proposition 4.1** Let f be any reduced closed strand diagram. Then:

- 1. Every component of f has at least one directed cycle.
- 2. Every directed cycle in f is either a free loop, a split loop, or a merge loop.
- 3. Any two directed cycles in f are disjoint, and no directed cycle intersects itself.

*Proof* For statement (1), observe that every vertex in  $\mathfrak f$  has at least one outgoing edge. Therefore, there exists an infinite directed path beginning at each vertex of  $\mathfrak f$ . Such a directed path must eventually intersect itself, and therefore  $\mathfrak f$  must contain a directed cycle.

For statement (2), suppose to the contrary that some directed cycle of f has both merges and splits. Then at least one edge of the cycle must begin at a merge and end at a split. Then this edge is subject to a type II reduction, which contradicts the assumption that f is reduced.

For statement (3), observe that any two intersecting directed cycles would have to merge together and then split apart, again contradicting the assumption that  $\mathfrak{f}$  is reduced. The same reasoning shows that a directed cycle cannot intersect itself.

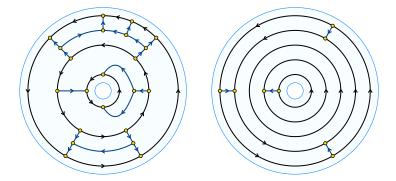


Fig. 4.1: A pair of reduced annular strand diagrams. The first is connected, while the second has three connected components. The directed cycles are shown in black.

Thus every reduced closed strand diagram consists of finitely many disjoint directed cycles, joined together by acyclic directed subgraphs. A bit more can be said for reduced annular and toral strand diagrams:

- For a reduced annular strand diagram, every component contains a directed cycle, and
  therefore every component must surround the central hole. It follows that every annular
  strand diagram consists of finitely many concentric components, each of which is itself
  an annular strand diagram. Note that a single component may have arbitrarily many
  directed cycles, though in a reduced diagram these cycles must alternate concentrically
  between merge loops and split loops.
- 2. For a reduced toral strand diagram, we know that every directed cycle must have positive winding number around the central hole. Since the directed cycles cannot intersect, all the directed cycles in a reduced toral strand diagram must be homotopic as loops in the torus. Specifically, each directed cycle must rotate around the central hole n times, and must rotate around the torus k times in the other direction, where k and n are relatively prime. By performing Dehn twists, we may assume that  $0 \le k < n$ . This gives us a rational invariant  $k/n \in \mathbb{Q} \cap [0,1)$  of the conjugacy class, which we refer to as the **rotation number**. Though we shall not prove it here, the rotation number of a conjugacy class for T is the same as the dynamical rotation number of any element of the class, viewed as a homeomorphism of the circle.

#### 5 Dynamics in Thompson's group F

In this section we show how the structure of an annular strand diagram is related to the dynamics of elements of the corresponding conjugacy class. For simplicity, we restrict to the case of Thompson's group F, although similar results hold for T and V.

#### 5.1 Strand diagrams as piecewise-linear functions

We begin by describing the relationship between strand diagrams and piecewise-linear homeomorphisms for elements of Thompson's group F. Given a strand diagram for an element  $f \in F$  and a real number  $t \in [0,1]$ , we can compute the image f(t) using the procedure shown in Figure 5.1(a). Roughly speaking, the strand diagram acts like a computer circuit: whenever a number  $t \in [0,1]$  is entered into the top, the signal winds its way through the circuit, emerging from the bottom as f(t). The path that the signal takes is determined by its binary digits, which change as the signal passes through each node, as shown in Figure 5.1(b).

Figure 5.2 shows the three different paths that a signal might take through the strand diagram for a certain element of F. Each of these paths corresponds to the linear mapping of an interval of the domain subdivision to the corresponding interval of the range subdivision.

Note that reductions do not change the action of the strand diagram on binary sequences, as shown in Figure 5.3. Thus equivalent strand diagrams really do represent the same piecewise-linear homeomorphism.

*Note* 5.1 More generally, we can interpret an (m,n)-strand diagram as a **Thompson-like homeomorphism**  $[0,m] \to [0,n]$ , i.e. a piecewise-linear homeomorphism whose slopes are powers of 2, and whose breakpoints have dyadic rational coordinates. Each number of the form  $(k-1)+(\alpha)$  corresponds to an input of  $\alpha$  entered into the kth source, or an output of  $\alpha$  emerging from the kth sink.

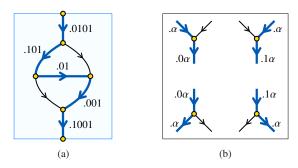


Fig. 5.1: (a) Computing f(.0101) using a strand diagram. (b) Rules for merges and splits.

# 5.2 Fixed points and replacement rules

Figure 5.4 shows a typical element of Thompson's group F. This element has four fixed points at 0, 1, 1/3, and 3/4, whose dynamics we now consider:

1. The fixed point at 0 is attracting, since the slope is 1/2. Near the fixed point, the function acts on binary digits via the rule  $f(.\alpha) = .0\alpha$ . This causes orbits to converge to 0:

$$.\alpha \ \mapsto \ .0\alpha \ \mapsto \ .00\alpha \ \mapsto \ .000\alpha \ \mapsto \ \cdots$$

The fixed point at 1 is similar, with  $f(.\alpha) = .1\alpha$  in a neighborhood of 1.

2. The fixed point at  $1/3 = .\overline{10}...$  is repelling, since the slope is 4. Near the fixed point, the function acts on binary digits via the rule  $f(.10\alpha) = .\alpha$ , which causes points to move away from 1/3:

$$.101010\alpha \ \mapsto \ .1010\alpha \ \mapsto \ .10\alpha \ \mapsto \ .\alpha \ \mapsto \ \cdots$$

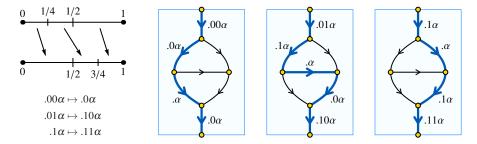


Fig. 5.2: Three paths through a strand diagram.

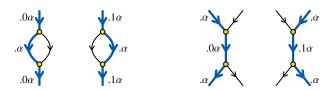


Fig. 5.3: Reductions do not change the underlying map.

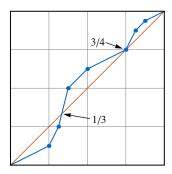


Fig. 5.4: An element of F.

3. Since 3/4 is a dyadic fraction, it has two binary expansions, namely  $.10\overline{1}$  and  $.11\overline{0}$ . Since the function has a breakpoint at 3/4, the behavior is different on the two sides of the fixed point. On the left side, the fixed point is attracting, with  $f(.10\alpha) = .101\alpha$ :

$$.10\alpha \mapsto .101\alpha \mapsto .1011\alpha \mapsto .10111\alpha \mapsto \cdots$$

On the right side, the fixed point is repelling, with  $f(.1100\alpha) = .110\alpha$ :

$$.110000\alpha \mapsto .11000\alpha \mapsto .1100\alpha \mapsto .110\alpha \mapsto \cdots$$

In general, an isolated fixed point for an element of F is either repelling or attracting, though it may have two different behaviors if the fixed point is dyadic. This behavior is evident from the action on binary digits: a repelling fixed point adds binary digits, while an attracting fixed point removes them.

Note that it is also possible for an element of *F* to have one or more closed intervals of fixed points. We refer to a maximal interval of this form as a **fixed interval**. Note that the endpoints of a fixed interval must always be dyadic.

It turns out that the behavior of the fixed points for an element  $f \in F$  is invariant under conjugacy. Indeed, it is possible to read this behavior directly from the annular strand diagram:

**Theorem 5.2** *There is a correspondence between the fixed points of an element*  $f \in F$  *and the closed loops in the corresponding reduced annular strand diagram. In particular:* 

- 1. Every fixed interval corresponds to a free loop in the annular strand diagram.
- 2. Every isolated non-dyadic fixed point (as well as 0 and 1) corresponds to either a split loop or a merge loop. In particular, a repelling fixed point with slope  $2^n$  corresponds to a split loop with n splits, and an attracting fixed point with slope  $2^{-n}$  corresponds to a merge loop with n merges.
- 3. Finally, every isolated dyadic fixed point (other than 0 or 1) corresponds to a pair of concentric loops, which may be split loops, merge loops, or one of each, depending on the behavior on the two sides of the fixed point.

In the latter two cases, the pattern of outward and inward connections around the loop determines the tail of the binary expansion of the fixed point. Specifically, each outward connection corresponds to a 1, and each inward connection corresponds to a 0.

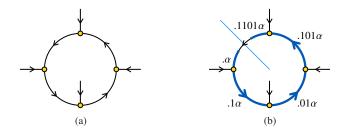


Fig. 5.5: (a) An example of a merge loop (b) Traveling around the merge loop

**Proof** By Proposition 3.2, there exists an (n,n)-strand diagram  $\mathfrak g$  whose closure is the reduced annular strand diagram for f. This diagram  $\mathfrak g$  corresponds to a Thompson-like homeomorphism  $g \colon [0,n] \to [0,n]$  which is conjugate to f. The fixed points of g are in one-to-one correspondence with the fixed points of f, and indeed it suffices to prove the statement of the theorem for the homeomorphism g.

To begin, suppose that the reduced annular strand diagram for g contains a merge loop, e.g. the loop shown in Figure 5.5(a). In the strand diagram for g, this loop corresponds to a directed path from the kth source to the kth sink, all of whose vertices are merges. In particular, if we feed a binary number  $\alpha$  into the kth source, it will emerge from the kth sink with some finite prefix added. For example, Figure 5.5(b) shows how an input of  $\alpha$  into the kth source will become output of .1101 $\alpha$  from the kth sink. In terms of functions, this means that

$$g((k-1)+.\alpha) = (k-1)+(.1101\alpha)$$

for all  $\alpha \in [0, 1]$ . It follows that  $(k-1) + \overline{1101}$  is an attracting fixed point for g.

In this way, each merge loop corresponds to an attracting fixed point of g, and each split loop corresponds to a repelling fixed point for g. A free loop corresponds to an entire unit interval [k,k+1] of fixed points. Note that every fixed point arises in this fashion. In particular, every fixed point must lie in some interval [k,k+1], and such an interval contains a fixed point if and only if there is a directed path in the strand diagram from the kth source to the kth sink. Also, note that each interval [k,k+1] contains at most one fixed point, and such a fixed point is dyadic if and only if it is an endpoint of the interval.

Finally, observe that the conjugacy between f and g preserves which fixed points are dyadic, and also preserves the tail of the binary expansion of each fixed point.

Note that the outermost loop of an annular strand diagram for an element of F corresponds to the fixed point  $0=.0000\cdots$ , while the innermost loop corresponds to the fixed point  $1=.1111\cdots$ . Within each connected component, the outermost and innermost loops correspond to dyadic fixed points, while the interior loops correspond to non-dyadic fixed points.

**Corollary 5.3** *Let*  $\mathfrak{f}'$  *be the reduced annular strand diagram for an element*  $f \in F$ . *Then every component of*  $\mathfrak{f}'$  *corresponds to exactly one of the following:* 

- 1. A maximal closed interval of fixed points of f (for a free loop), or
- 2. A maximal interval with no dyadic fixed points of f in its interior.

If  $f \in F$ , a **cut point** of f is either an isolated dyadic fixed point of f, or an endpoint of a maximal interval of fixed points. If  $0 = \alpha_0 < \alpha_1 < \dots < \alpha_n = 1$  are the cut points of f, then the restrictions  $f_i \colon [\alpha_{i-1}, \alpha_i] \to [\alpha_{i-1}, \alpha_i]$  are called the **components** of f (see Figure 5.6). Each component of f corresponds to one connected component of the reduced annular strand diagram (figure 5.6). If  $\alpha < \beta$  are any dyadic rationals, it is well known (see [7]) that there exists a Thompson-like homeomorphism  $\varphi \colon [\alpha, \beta] \to [0, 1]$ . It follows that any Thompson-like homeomorphism of  $[\alpha, \beta]$  can be conjugated by  $\varphi$  to give an element of F. The following are straightforward.

**Proposition 5.4** Let  $f \in F$  have components  $f_i \colon [\alpha_{i-1}, \alpha_i] \to [\alpha_{i-1}, \alpha_i]$ , and let f' be the reduced annular strand diagram for f. Then for each i, the component of S corresponding to  $f_i$  is the reduced annular strand diagram for any element of F conjugate to  $f_i$ .

**Corollary 5.5** Let  $f,g \in F$  have components  $f_1, \ldots, f_n$  and  $g_1, \ldots, g_n$ . Then f is conjugate to g in F if and only if each  $f_i$  is conjugate to  $g_i$  through some Thompson-like homeomorphism.

Remark 5.6 Although the work outlined in Sections 5.1 and 5.2 describes a relation between the dynamics of elements of F and annular strand diagrams, one can describe a similar connection between elements in T (respectively, V) and toral strand diagrams (respectively, abstract closed strand diagrams). Bleak  $et\ al.$  give a description of centralizers for Thompson's group V based on a study of the orbits of an element [2]. The description in [2] recovers abstract closed strand diagrams and shows how they encode information about the dynamics of all the elements of a conjugacy class, giving a generalization of the results obtained above for F.

### 5.3 Mather Invariants

Conjugacy in F was first investigated by Brin and Squier [5], who successfully found an invariant for conjugacy in the full group of piecewise-linear homeomorphisms of the interval with finitely many breakpoints. Their invariant was based on some ideas of Mather [14] for

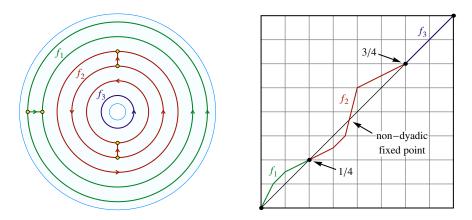


Fig. 5.6: An element of F with three components, and the corresponding annular strand diagram.

determining whether two given diffeomorphisms of the unit interval are conjugate. In this section we show that the annular strand diagrams described in Section 2 can be used to define a Mather-type invariant for elements of F for functions without fixed points (except for 0 and 1). The invariant that we describe is a suitable adaptation of Brin and Squier which works for F. independently from our work, Gill and Short [8] also successfully found a way to generalize Brin and Squier's version of the Mather invariant and thus providing a description of conjugacy in F which is similar to the one presented in the current section (although the proofs are somewhat different).

Recall that a map  $f \in F$  is called a **one-bump function** is an element such that f(x) > x for all  $x \in (0,1)$ . Consider a one-bump function  $f \in F$ , with slope  $2^{m_0}$  at 0 and slope  $2^{m_1}$  at 1. In a neighborhood of zero, f acts as multiplication by  $2^{m_0}$ ; in particular, for any sufficiently small t > 0, the interval  $[t, 2^{m_0}t]$  is a fundamental domain for the action of f. If we make the identification  $t \sim 2^{m_0}t$  in the interval  $(0, \varepsilon)$ , we obtain a circle  $C_0$ , with partial covering map  $p_0 \colon (0, \varepsilon) \to C_0$ . Similarly, if we identify  $(1 - t) \sim (1 - 2^{m_1}t)$  on the interval  $(1 - \delta, 1)$ , we obtain a circle  $C_1$ , with partial covering map  $p_1 \colon (1 - \delta, 1) \to C_1$ .

If N is sufficiently large, then  $f^N$  will take some lift of  $C_0$  to  $(0, \varepsilon)$  and map it to the interval  $(1 - \delta, 1)$ . This induces a map  $f^{\infty}$ :  $C_0 \to C_1$ , making the following diagram commute:

$$(0,\varepsilon) \xrightarrow{f^N} (1-\delta,1)$$

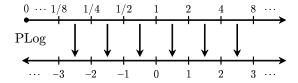
$$p_0 \downarrow \qquad \qquad \downarrow p_1$$

$$C_0 \xrightarrow{f^\infty} C_1$$

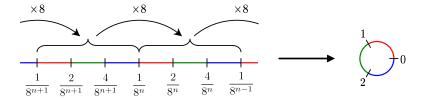
We note that  $f^{\infty}$  does not depend on the specific value of N chosen. Any map  $f^s$ , for  $s \ge N$ , induces the same map  $f^{\infty}$ . This is because f acts as the identity on  $C_1$  by construction and  $f^s$  can be written as  $f^{s-N}(f^N(t))$ , with  $f^N(t) \in (1-\delta,1)$ .

Remark 5.7 The map  $f^{\infty}$  defined above is the **Mather invariant** for f as defined in Brin and Squier [5] for functions in the full group of piecewise-linear homeomorphisms of the interval with finitely many breakpoints. Since f is an element of F, we will now rescale the two circles  $C_0$  and  $C_1$  so that their length is equal to the exponents of the slopes at 0 and 1.

**Definition 5.8** The **piecewise-linear logarithm** PLog:  $(0,\infty) \to (-\infty,\infty)$  is the piecewise-linear function that maps the interval  $[2^k,2^{k+1}]$  linearly onto [k,k+1] for every  $k \in \mathbb{Z}$ , as shown below:



Suppose that  $f \in F$  is a one-bump function with slope  $2^m$  at 0 and slope  $2^{-n}$  at 1, and let  $f^{\infty} \colon C_0 \to C_1$  be the corresponding Mather invariant. In a neighborhood of 0, the function f acts as multiplication by  $2^m$ . In particular, PLog f(t) = m + PLog f for all  $t \in (0, \varepsilon)$ , so we can identify  $C_0$  with the circle  $\mathbb{R}/m\mathbb{Z}$ . The following picture shows the case m = 3:



In a similar way, we can use the function  $t \mapsto -\operatorname{PLog}(1-t)$  to identify  $C_1$  with the circle  $\mathbb{R}/n\mathbb{Z}$ . This lets us regard the Mather invariant for f as a function  $f^{\infty} \colon \mathbb{R}/m\mathbb{Z} \to \mathbb{R}/n\mathbb{Z}$ . Because  $f^N$  and PLog are piecewise-linear, the Mather invariant  $f^{\infty}$  is a piecewise-linear function. Moreover,  $f^{\infty}$  is Thompson-like: all the slopes are powers of 2, and the breakpoints are dyadic rational numbers of  $\mathbb{R}/m\mathbb{Z} = [0,m]/\{0,m\}$ .

Now, if  $k \in \mathbb{Z}$ , then the map  $t \mapsto 2^k t$  on  $(0, \varepsilon)$  induces precisely an integer rotation of  $\mathbb{R}/m\mathbb{Z}$ :

$$rot_k(\theta) = \theta + k \mod m$$

We are now ready to state a criterion for conjugacy via Mather invariants:

**Theorem 5.9** Let  $f,g \in F$  be one-bump functions with  $f'(0) = g'(0) = 2^m$  and  $f'(1) = g'(1) = 2^{-n}$ , and let  $f^{\infty}, g^{\infty} \colon \mathbb{R}/m\mathbb{Z} \to \mathbb{R}/n\mathbb{Z}$  be the corresponding Mather invariants. Then f and g are conjugate if and only if  $f^{\infty}$  and  $g^{\infty}$  differ by integer rotations of the domain and range circles:

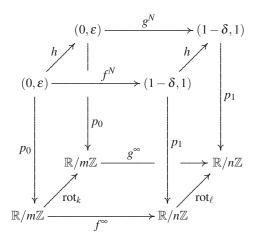
$$\mathbb{R}/m\mathbb{Z} \xrightarrow{f^{\infty}} \mathbb{R}/n\mathbb{Z}$$

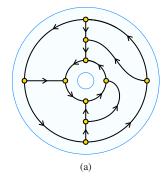
$$\operatorname{rot}_{k} \downarrow \qquad \qquad \downarrow \operatorname{rot}_{\ell}$$

$$\mathbb{R}/m\mathbb{Z} \xrightarrow{g^{\infty}} \mathbb{R}/n\mathbb{Z}$$

We only sketch the proof of Theorem 5.9 as some steps are straightforward, while others can be obtained using small variations of some of our previous arguments.

The forward direction follows observing that, if  $f = h^{-1}gh$  for some  $h \in F$ , then the following diagram commutes, where  $k = \log_2 h'(0)$  and  $\ell = \log_2 h'(1)$ :





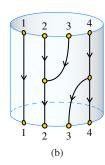


Fig. 5.7: (a) Annular strand diagram for a one-bump function (b) The corresponding cylindrical strand diagram

For the converse, we must show that any two one-bump functions whose Mather invariants differ by integer rotation are conjugate in F. To prove this, we describe an explicit correspondence between Mather invariants and reduced annular strand diagrams.

If  $f \in F$  is a one-bump function, then the only fixed points of f are at 0 and 1. Therefore, the reduced annular strand diagram for f has only two directed cycles (see Figure 5.7a). Since f'(0) > 1, the outer cycle (corresponding to 0) must be a split loop, and the inner cycle (corresponding to 1) must be a merge loop. If we remove these two cycles, we get an (m,n)-cylindrical strand diagram (see Figure 5.7b). Such a diagram can be used to describe a Thompson-like map between two circles. The following result only requires one to observe how to construct and glue forests to obtain a cylindrical strand diagram and conversely how to cut such diagram to obtain a pair of forests. We omit its proof.

#### **Proposition 5.10** There is a one-to-one correspondence between

- 1. Reduced cylindrical (m,n)-strand diagrams, and
- 2. Thompson-like functions  $\mathbb{R}/m\mathbb{Z} \to \mathbb{R}/n\mathbb{Z}$ , with two functions considered equivalent if they differ by integer rotation of the domain and range circles.

To complete the proof of Theorem 5.9 we need the following result which can be obtained by imitating the idea of Theorem 5.2 to study the image of a binary number under many iterations.

**Proposition 5.11** Let  $\mathscr{A}$  be the reduced annular strand diagram for a one-bump function  $f \in F$ , and let  $\mathscr{C}$  be the cylindrical (m,n)-strand diagram obtained by removing the merge and split loops from  $\mathscr{A}$ . Then  $\mathscr{C}$  is the cylindrical strand diagram for the Mather invariant  $f^{\infty} \colon \mathbb{R}/m\mathbb{Z} \to \mathbb{R}/n\mathbb{Z}$ .

Remark 5.12 Brin and Squier used the Mather invariant as a means to describe a conjugacy invariant in the full group of piecewise-linear homeomorphisms of the unit interval with finitely many breakpoints. With Theorem 5.9, the authors found a way to adapt Brin and Squier's invariant to the case of Thompson's group F. After the authors found the description of Theorem 5.9 and proved it using strand diagrams, the second author analyzed this description using the techniques appearing in the joint paper of Kassabov and the second author [13]: this resulted in a generalization [16] which works for a larger class of groups of piecewise-linear homeomorphisms (essentially the Thompson-Stein groups) and bridging the gap between the descriptions in [5] and Theorem 5.9.

#### Acknowledgments

The authors would like to thank Ken Brown and Martin Kassabov for a very careful reading of this paper and many helpful remarks. The authors would also like to thank Stephen Pride for providing fundamental references and Collin Bleak, Matt Brin, Jörg Lehnert and Mark Sapir for helpful conversations. Finally, the authors would like to thank the referee for many helpful comments and suggestions.

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